

# Plan for today

- Dynamic Programming
  - Employment Scheduling Problem
  - Reject Allowance Problem
  - Winning in Vegas Problem

# DYNAMIC PROGRAMMING

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... where we see the algorithmic technique of dynamic programming (through examples)

# EMPLOYMENT SCHEDULING



... a problem where the number of states is infinity and the decision variables can take real values

# Employment Scheduling Problem

- Consider a local shop that is subject to seasonal fluctuation
- At the end of Spring, the shop has 255 workers

Required number of workers in each season

Season	Summer	Fall	Winter	Spring
Requirements	220	240	200	255

- Requirement has to be met in all seasons,
  - i.e., cannot hire below the numbers given in the above table, but can hire more than the requirement
- The planning horizon ends next spring; should have exactly 255 workers for Spring
- Costs:
  - Hiring or firing workers costs  $2x^2$ , where  $x$  is the number of workers hired/fired
  - Each worker in excess of the requirement costs \$20 per season
- Fractional employment levels are allowed (due to part-time employees)
  - cost also applies on a fractional basis

**Question: determine the number of workers to hire/fire at the beginning of each season to minimize the total cost**

# Stages, States, Decision vars, Optimality criterion

- Stages: four stages, representing decision-making at the beginning of summer, fall, winter, and spring
- States:  $s_n$  represents the number of workers at the beginning of  $n$ 'th stage

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 = 255$	$x_1$	$s_1 + x_1$	<b>220</b>
2. Fall	$s_2 = s_1 + x_1$	$x_2$	$s_2 + x_2$	<b>240</b>
3. Winter	$s_3 = s_2 + x_2$	$x_3$	$s_3 + x_3$	<b>200</b>
4. Spring	$s_4 = s_3 + x_3$	$x_4$	$s_4 + x_4 = 255$	<b>255</b>

- Decision Variables/Policy:  $x_n =$  Number of workers hired/fired at stage  $n$
- Optimality criterion: minimize total cost

## Value function

Need:  $f_1^*(255)$ 

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 (=255)$	$x_1$	$s_1+x_1$	<b>220</b>
2. Fall	$s_2 (=s_1+x_1)$	$x_2$	$s_2+x_2$	<b>240</b>
3. Winter	$s_3 (=s_2+x_2)$	$x_3$	$s_3+x_3$	<b>200</b>
4. Spring	$s_4 (=s_3+x_3)$	$x_4$	$s_4+x_4=255$	<b>255</b>

$f_n(s_n, x_n)$  = Cost of starting from state  $s_n$  in stage  $n$ , hiring/firing  $x_n$  in stage  $n$  and making optimum hiring/firing decisions thereafter

$$f_n^*(s_n) = \min_{x_n: s_n+x_n \geq \text{requirement at stage } n} f_n(s_n, x_n)$$

Need:  $f_1^*(255)$

## Value function: final stage

Need:  $f_1^*(255)$ 

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 (=255)$	$x_1$	$s_1+x_1$	<b>220</b>
2. Fall	$s_2 (=s_1+x_1)$	$x_2$	$s_2+x_2$	<b>240</b>
3. Winter	$s_3 (=s_2+x_2)$	$x_3$	$s_3+x_3$	<b>200</b>
4. Spring	$s_4 (=s_3+x_3)$	$x_4$	$s_4+x_4=255$	<b>255</b>

$f_n(s_n, x_n)$  = Cost of starting from state  $s_n$  in stage  $n$ , hiring/firing  $x_n$  in stage  $n$  and making optimum hiring/firing decisions thereafter

$$f_n^*(s_n) = \min_{x_n: s_n+x_n \geq \text{requirement at stage } n} f_n(s_n, x_n)$$

Need:  $f_1^*(255)$ 

$$f_4^*(s_4) = 2x_4^2 = 2(255 - s_4)^2$$

## Value function: recursive relationship

Need:  $f_1^*(255)$ 

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 (=255)$	$x_1$	$s_1+x_1$	<b>220</b>
2. Fall	$s_2 (=s_1+x_1)$	$x_2$	$s_2+x_2$	<b>240</b>
3. Winter	$s_3 (=s_2+x_2)$	$x_3$	$s_3+x_3$	<b>200</b>

$f_n(s_n, x_n)$  = Cost of starting from state  $s_n$  in stage  $n$ , hiring/firing  $x_n$  in stage  $n$  and making optimum hiring/firing decisions thereafter

$$f_n^*(s_n) = \min_{x_n: s_n+x_n \geq \text{requirement at stage } n} f_n(s_n, x_n)$$

Need:  $f_1^*(255)$ 

$$f_3(s_3, x_3) = \underbrace{2x_3^2 + 20(s_3 + x_3 - 200)}_{\text{Immediate cost of starting at } s_3 \text{ and hiring/firing } x_3 \text{ in stage 3}} + f_4^*(s_3 + x_3)$$

$$f_3^*(s_3) = \min_{x_3: s_3+x_3 \geq 200} \{f_3(s_3, x_3)\}$$

$f_3(s_3, x_3) =$   
 Immediate cost of starting at  $s_3$   
 and hiring/firing  $x_3$  in stage 3  
 +  
 the cost of optimal decision if we  
 start stage 4 from state  $s_3 + x_3$



## Value function: recursive relationship

Need:  $f_1^*(255)$ 

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 (=255)$	$x_1$	$s_1+x_1$	<b>220</b>
2. Fall	$s_2 (=s_1+x_1)$	$x_2$	$s_2+x_2$	<b>240</b>
3. Winter	$s_3 (=s_2+x_2)$	$x_3$	$s_3+x_3$	<b>200</b>

$f_n(s_n, x_n)$  = Cost of starting from state  $s_n$  in stage  $n$ , hiring/firing  $x_n$  in stage  $n$  and making optimum hiring/firing decisions thereafter

$$f_n^*(s_n) = \min_{x_n: s_n+x_n \geq \text{requirement at stage } n} f_n(s_n, x_n)$$

Need:  $f_1^*(255)$ 

$$f_3(s_3, x_3) = 2x_3^2 + 20(s_3 + x_3 - 200) + f_4^*(s_3 + x_3)$$

$$f_3^*(s_3) = \min_{x_3: s_3+x_3 \geq 200} \{f_3(s_3, x_3)\}$$

$$f_n(s_n, x_n) = 2x_n^2 + 20(s_n + x_n - \text{requirement at stage } n) + f_{n+1}^*(s_n + x_n)$$

$$f_n^*(s_n) = \min_{x_n: s_n+x_n \geq \text{requirement at stage } n} \{f_n(s_n, x_n)\}$$

## Solution procedure: backward induction (stage 3)

Need:  $f_1^*(255)$ 

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 (=255)$	$x_1$	$s_1+x_1$	<b>220</b>
2. Fall	$s_2 (=s_1+x_1)$	$x_2$	$s_2+x_2$	<b>240</b>
3. Winter	$s_3 (=s_2+x_2)$	$x_3$	$s_3+x_3$	<b>200</b>

Given that  $f_4^*(s_4) = 2x_4^2 = 2(255 - s_4)^2$

$$f_3(s_3, x_3) = 2x_3^2 + 20(s_3 + x_3 - 200) + f_4^*(s_3 + x_3)$$

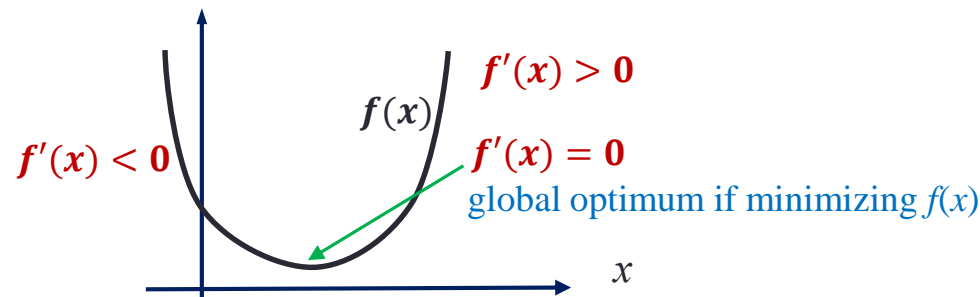
$$\begin{aligned} f_3^*(s_3) &= \min_{x_3: s_3+x_3 \geq 200} \{2x_3^2 + 20(s_3 + x_3 - 200) + f_4^*(s_3 + x_3)\} \\ &= \min_{x_3: s_3+x_3 \geq 200} \{2x_3^2 + 20(s_3 + x_3 - 200) + 2(255 - s_3 - x_3)^2\} \end{aligned}$$

Ignoring the  $s_3 + x_3 \geq 200$  constraint, the optimum value is attained at  $x_3^* =$



# Minimizing Single Variable Functions

- Continuous function  $f(x)$



1. Find a point  $x^*$  where  $f'(x^*) = 0$
2. Verify if  $x^*$  is a minimizer (or a maximizer)  
... by checking if  $f''(x^*) > 0$  (or  $f''(x^*) < 0$ )

## Solution procedure: backward induction (stage 3)

Need:  $f_1^*(255)$ 

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 (=255)$	$x_1$	$s_1+x_1$	<b>220</b>
2. Fall	$s_2 (=s_1+x_1)$	$x_2$	$s_2+x_2$	<b>240</b>
3. Winter	$s_3 (=s_2+x_2)$	$x_3$	$s_3+x_3$	<b>200</b>

Given that  $f_4^*(s_4) = 2x_4^2 = 2(255 - s_4)^2$

$$f_3(s_3, x_3) = 2x_3^2 + 20(s_3 + x_3 - 200) + f_4^*(s_3 + x_3)$$

$$\begin{aligned} f_3^*(s_3) &= \min_{x_3: s_3+x_3 \geq 200} \{2x_3^2 + 20(s_3 + x_3 - 200) + f_4^*(s_3 + x_3)\} \\ &= \min_{x_3: s_3+x_3 \geq 200} \{2x_3^2 + 20(s_3 + x_3 - 200) + 2(255 - s_3 - x_3)^2\} \end{aligned}$$

Ignoring the  $s_3 + x_3 \geq 200$  constraint, the optimum value is attained at  $x_3^* =$

$$g(x_3) := f_3(s_3, x_3) = 2x_3^2 + 20(s_3 + x_3 - 200) + 2(255 - s_3 - x_3)^2$$

$$\begin{aligned} \frac{dg(x_3)}{dx_3} &= 4x_3 + 20 - 4(255 - s_3 - x_3) = 8x_3 + 4s_3 - 1040 \\ &= 8 \left( x_3 - 125 + \frac{1}{2}s_3 \right) \end{aligned}$$

$$\frac{d^2g(x_3)}{dx_3^2} = 8 > 0 \Rightarrow g \text{ is minimized at } x_3 = 125 - \frac{1}{2}s_3$$



## Solution procedure: backward induction (stage 3)

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 (=255)$	$x_1$	$s_1+x_1$	<b>220</b>
2. Fall	$s_2 (=s_1+x_1)$	$x_2$	$s_2+x_2$	<b>240</b>
3. Winter	$s_3 (=s_2+x_2)$	$x_3$	$s_3+x_3$	<b>200</b>

Given that  $f_4^*(s_4) = 2x_4^2 = 2(255 - s_4)^2$

$$f_3(s_3, x_3) = 2x_3^2 + 20(s_3 + x_3 - 200) + f_4^*(s_3 + x_3)$$

$$\begin{aligned} f_3^*(s_3) &= \min_{x_3: s_3+x_3 \geq 200} \{ 2x_3^2 + 20(s_3 + x_3 - 200) + f_4^*(s_3 + x_3) \} \\ &= \min_{x_3: s_3+x_3 \geq 200} \{ 2x_3^2 + 20(s_3 + x_3 - 200) + 2(255 - s_3 - x_3)^2 \} \end{aligned}$$

Ignoring the  $s_3 + x_3 \geq 200$  constraint, the optimum value is attained at  $x_3^* = 125 - \frac{1}{2}s_3$

Constraint:  $s_3 + x_3 \geq 200$ ,

Substituting  $x_3^* = 125 - \frac{1}{2}s_3$ , we need  $125 + \frac{1}{2}s_3 \geq 200$ , i.e., we need  $s_3 \geq 150$ .

Note that this is automatically satisfied since  $s_3 \geq 240$ !

So we plug the solution for  $x_3^*$  into the value function

$$f_3^*(s_3) = 2 \left( 125 - \frac{1}{2}s_3 \right)^2 + 20 \left( \frac{1}{2}s_3 - 75 \right) + 2 \left( 130 - \frac{1}{2}s_3 \right)^2$$



## Solution procedure: backward induction (stage 2)

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 (=255)$	$x_1$	$s_1+x_1$	<b>220</b>
2. Fall	$s_2 (=s_1+x_1)$	$x_2$	$s_2+x_2$	<b>240</b>

Given that  $f_3^*(s_3) = 2 \left(125 - \frac{1}{2}s_3\right)^2 + 20 \left(\frac{1}{2}s_3 - 75\right) + 2 \left(130 - \frac{1}{2}s_3\right)^2$

$$f_2(s_2, x_2) = 2x_2^2 + 20(s_2 + x_2 - 240) + f_3^*(s_2 + x_2)$$

$$= 2x_2^2 + 20(s_2 + x_2 - 240) + 2 \left(125 - \frac{1}{2}(s_2 + x_2)\right)^2 + 20 \left(\frac{1}{2}(s_2 + x_2) - 75\right) + 2 \left(130 - \frac{1}{2}(s_2 + x_2)\right)^2$$

$$f_2^*(s_2) = \min_{x_2: s_2+x_2 \geq 240} \{f_2(s_2, x_2)\}$$

$$g(x_2) := f_2(s_2, x_2)$$

$$g(x_2) = 2x_2^2 + 20(s_2 + x_2 - 240) + 2 \left(125 - \frac{1}{2}(s_2 + x_2)\right)^2 + 20 \left(\frac{1}{2}(s_2 + x_2) - 75\right) + 2 \left(130 - \frac{1}{2}(s_2 + x_2)\right)^2$$

$$\frac{dg(x_2)}{dx_2} = 4x_2 + 20 + (-2) \left(125 - \frac{1}{2}(s_2 + x_2)\right) + 10 + (-2) \left(130 - \frac{1}{2}(s_2 + x_2)\right)$$

$$\frac{dg(x_2)}{dx_2} = 6x_2 + 2s_2 - 480 \Rightarrow \frac{dg(x_2)}{dx_2} = 0 \Rightarrow x_2 = 80 - \frac{1}{3}s_2$$

$$\frac{d^2g(x_2)}{dx_2^2} = 6 > 0 \Rightarrow g(x_2) \text{ is minimized at } x_2 = 80 - \frac{1}{3}s_2 \Rightarrow x_2^* = 80 - \frac{1}{3}s_2 \text{ if } x_2^* + s_2 \geq 240, \text{ i.e., if } s_2 \geq 240$$

Need:  $f_1^*(255)$

## Solution procedure: backward induction (stage 2)

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 (=255)$	$x_1$	$s_1+x_1$	<b>220</b>
2. Fall	$s_2 (=s_1+x_1)$	$x_2$	$s_2+x_2$	<b>240</b>

Given that  $f_3^*(s_3) = 2 \left(125 - \frac{1}{2}s_3\right)^2 + 20 \left(\frac{1}{2}s_3 - 75\right) + 2 \left(130 - \frac{1}{2}s_3\right)^2$

$$f_2(s_2, x_2) = 2x_2^2 + 20(s_2 + x_2 - 240) + f_3^*(s_2 + x_2)$$

$$= 2x_2^2 + 20(s_2 + x_2 - 240) + 2 \left(125 - \frac{1}{2}(s_2 + x_2)\right)^2 + 20 \left(\frac{1}{2}(s_2 + x_2) - 75\right) + 2 \left(130 - \frac{1}{2}(s_2 + x_2)\right)^2$$

$$f_2^*(s_2) = \min_{x_2: s_2+x_2 \geq 240} \{f_2(s_2, x_2)\}$$

If  $s_2 \leq 240$ , (we also need  $x_2 \geq 240 - s_2$ )

$$g(x_2) := f_2(s_2, x_2)$$

$$g(x_2) = 2x_2^2 + 20(s_2 + x_2 - 240) + 2 \left(125 - \frac{1}{2}(s_2 + x_2)\right)^2 + 20 \left(\frac{1}{2}(s_2 + x_2) - 75\right) + 2 \left(130 - \frac{1}{2}(s_2 + x_2)\right)^2$$

$$\frac{dg(x_2)}{dx_2} = 4x_2 + 20 + (-2) \left(125 - \frac{1}{2}(s_2 + x_2)\right) + 10 + (-2) \left(130 - \frac{1}{2}(s_2 + x_2)\right)$$

$$\frac{dg(x_2)}{dx_2} = 6x_2 + 2s_2 - 480$$

$$\geq 6(240 - s_2) + 2s_2 - 480 \text{ (since } x_2 \geq 240 - s_2)$$

$$= 960 - 4s_2 > 0 \text{ (since } s_2 \leq 240)$$

$\Rightarrow g(x_2)$  is increasing if we increase  $x_2$   
 $\Rightarrow g(x_2)$  is minimized at  $x_2 = 240 - s_2$

Need:  $f_1^*(255)$

# Solution procedure: backward induction (stage 2)

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 (=255)$	$x_1$	$s_1+x_1$	<b>220</b>
2. Fall	$s_2 (=s_1+x_1)$	$x_2$	$s_2+x_2$	<b>240</b>

Given that  $f_3^*(s_3) = 2 \left(125 - \frac{1}{2}s_3\right)^2 + 20 \left(\frac{1}{2}s_3 - 75\right) + 2 \left(130 - \frac{1}{2}s_3\right)^2$

$$f_2(s_2, x_2) = 2x_2^2 + 20(s_2 + x_2 - 240) + f_3^*(s_2 + x_2)$$

$$= 2x_2^2 + 20(s_2 + x_2 - 240) + 2 \left(125 - \frac{1}{2}(s_2 + x_2)\right)^2 + 20 \left(\frac{1}{2}(s_2 + x_2) - 75\right) + 2 \left(130 - \frac{1}{2}(s_2 + x_2)\right)^2$$

$$f_2^*(s_2) = \min_{x_2: s_2+x_2 \geq 240} \{f_2(s_2, x_2)\}$$

Solving the above minimization problem yields  $x_2^* = \begin{cases} 80 - \frac{s_2}{3} & \text{if } s_2 \geq 240 \\ 240 - s_2 & \text{if } s_2 < 240 \end{cases}$

Plugging  $x_2^*$  into  $f_2^*(s_2)$ , gives

$$f_2^*(s_2) = \begin{cases} 2 \left(80 - \frac{1}{3}s_2\right)^2 + 20(s_2 - 195) + 2 \left(85 - \frac{1}{3}s_2\right)^2 + 2 \left(90 - \frac{1}{3}s_2\right)^2 & \text{if } s_2 \geq 240 \\ 2(240 - s_2)^2 + 1150 & \text{if } s_2 < 240 \end{cases}$$



Need:  $f_1^*(255)$

## Solution procedure: backward induction (stage 1)

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 (=255)$	$x_1$	$s_1+x_1$	<b>220</b>

$$f_1(s_1 = 255, x_1) = 2x_1^2 + 20(255 + x_1 - 220) + f_2^*(255 + x_1)$$

$$f_1^*(s_1) = \min_{x_1: s_1+x_1 \geq 220} f_1(s_1, x_1)$$

Recall that the value function  $f_2^*(s_2)$  has two branches, depending on the value of  $s_2$ , so here we need to discuss two separate cases:

Case 1:  $s_2 < 240$

Case 2:  $s_2 \geq 240$

### Case 1:

If  $s_2 = s_1 + x_1 < 240$ , i.e.,  $x_1 < 240 - s_1 = -15$ , then

$$f_2^*(s_2) = f_2^*(s_1 + x_1) = 2(240 - 255 - x_1)^2 + 1150$$

Substituting, we get  $f_1(s_1, x_1) = 2x_1^2 + 20(255 + x_1 - 220) + 2(15 + x_1)^2 + 1150$

$$f_1^*(s_1) = \min_{x_1: s_1+x_1 \geq 220} \{f_1(s_1, x_1)\}$$

$$g(x_1) := f_1(s_1, x_1)$$

$$g(x_1) = 2x_1^2 + 20(35 + x_1 - 220) + 2(15 + x_1)^2 + 1150$$

$$\frac{dg(x_1)}{dx_1} = 4x_1 + 20 + 4(15 + x_1) = 8x_1 + 80$$

$$< 0 \text{ if } x_1 < -15$$

}  $\Rightarrow g(x_1)$  is decreasing if we increase  $x_1$   
 $\Rightarrow g(x_1)$  is minimized at  $x_1 = -15$ ,  
 i.e., in Case 2

## Solution procedure: backward induction (stage 1)

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 (=255)$	$x_1$	$s_1+x_1$	<b>220</b>

$$f_1(s_1 = 255, x_1) = 2x_1^2 + 20(255 + x_1 - 220) + f_2^*(255 + x_1)$$

$$f_1^*(s_1) = \min_{x_1: s_1+x_1 \geq 220} f_1(s_1, x_1)$$

### Case 2:

If  $s_2 = s_1 + x_1 \geq 240$ , i.e.,  $x_1 \geq -15$ , then

$$f_1(s_1, x_1) = 2x_1^2 + 20(35 + x_1) + 2\left(5 + \frac{1}{3}x_1\right)^2 + 20(60 + x_1) + 2\left(\frac{1}{3}x_1\right)^2 + 2\left(5 - \frac{1}{3}x_1\right)^2$$

$$f_1^*(s_1) = \min_{x_1: s_1+x_1 \geq 220} \{f_1(s_1, x_1)\}$$

$$g(x_1) := f_1(s_1, x_1)$$

$$g(x_1) = 2x_1^2 + 20(35 + x_1) + 2\left(5 + \frac{1}{3}x_1\right)^2 + 20(60 + x_1) + 2\left(\frac{1}{3}x_1\right)^2 + 2\left(5 - \frac{1}{3}x_1\right)^2$$

$$\frac{dg(x_1)}{dx_1} = 4x_1 + 20 + 4\left(5 + \frac{1}{3}x_1\right)\left(\frac{1}{3}\right) + 20 + 4\left(\frac{1}{3}x_1\right)\left(\frac{1}{3}\right) + 4\left(5 - \frac{1}{3}x_1\right)\left(-\frac{1}{3}\right)$$

$$\frac{dg(x_1)}{dx_1} = \frac{16}{3}x_1 + 40 \Rightarrow \frac{dg(x_1)}{dx_1} = 0 \Rightarrow x_1 = -\frac{15}{2}$$

$$\frac{d^2g(x_1)}{dx_1^2} = \frac{16}{3} > 0 \Rightarrow g(x_1) \text{ is minimized at } x_1 = -\frac{15}{2} \Rightarrow x_1^* = -\frac{15}{2} \text{ and } f_1^*(255) = 1850$$

# Optimal Policy

	beginning state	hire/fire	ending state	requirement
1. Summer	$s_1 (=255)$	$x_1$	$s_1+x_1$	<b>220</b>
2. Fall	$s_2 (=s_1+x_1)$	$x_2$	$s_2+x_2$	<b>240</b>
3. Winter	$s_3 (=s_2+x_2)$	$x_3$	$s_3+x_3$	<b>200</b>
4. Spring	$s_4 (=s_3+x_3)$	$x_4$	$s_4+x_4=255$	<b>255</b>

$$x_1^* = -7.5 \quad x_2^* = 80 - \frac{s_2}{3} \quad x_3^* = 125 - \frac{s_3}{2} \quad x_4^* = 255 - s_4$$

$$f_1^*(255) = 1850$$

Optimum cost is \$1,850

Optimum policy:

	beginning state	hire/fire	ending state	requirement	surplus
1. Summer	255	-7.5	247.5	220	27.5
2. Fall	247.5	-2.5	245	240	5
3. Winter	245	2.5	247.5	200	47.5
4. Spring	247.5	7.5	255	255	0

# PROBABILISTIC DP

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... problems where the transitions between states are probabilistic



# REJECT ALLOWANCE

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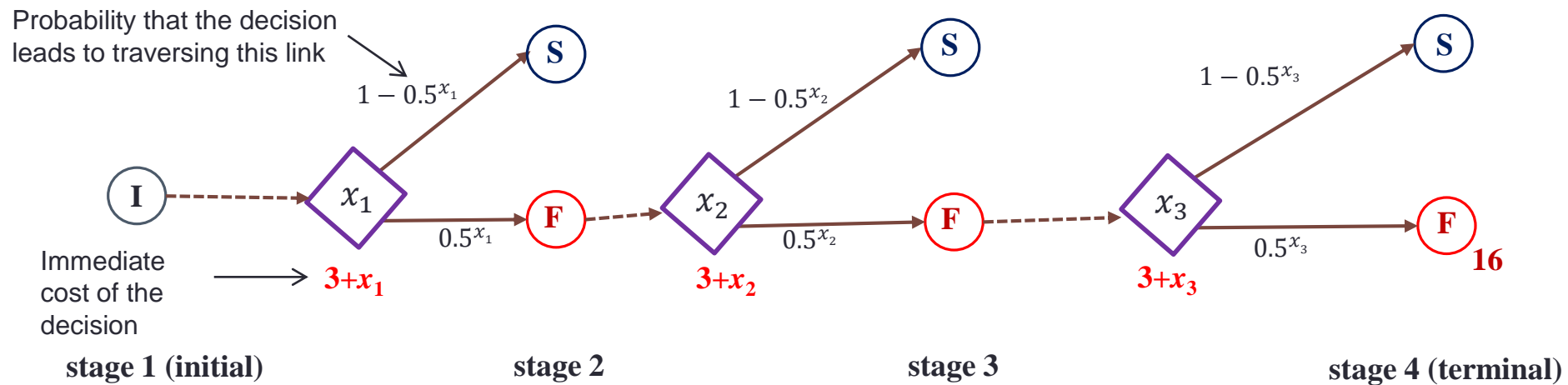
## Example 4: Reject allowance

- A customer orders one unit of the product with the option to reject it if it is defective
- To sell a single unit, the manufacturer produces a batch (containing several units of the product), hoping that at least one unit from the batch will NOT be defective
- Each produced unit will be defective with probability  $\frac{1}{2}$
- If the entire batch is defective, the manufacturer can run another round of production to produce another batch
- The maximum possible number of production runs is **three**
- Each production run costs **\$3** to set up and each unit costs **\$1** to produce
- If all three runs fail to produce a non-defective unit, the process ends and the manufacturer loses **\$16** of revenue from the customer

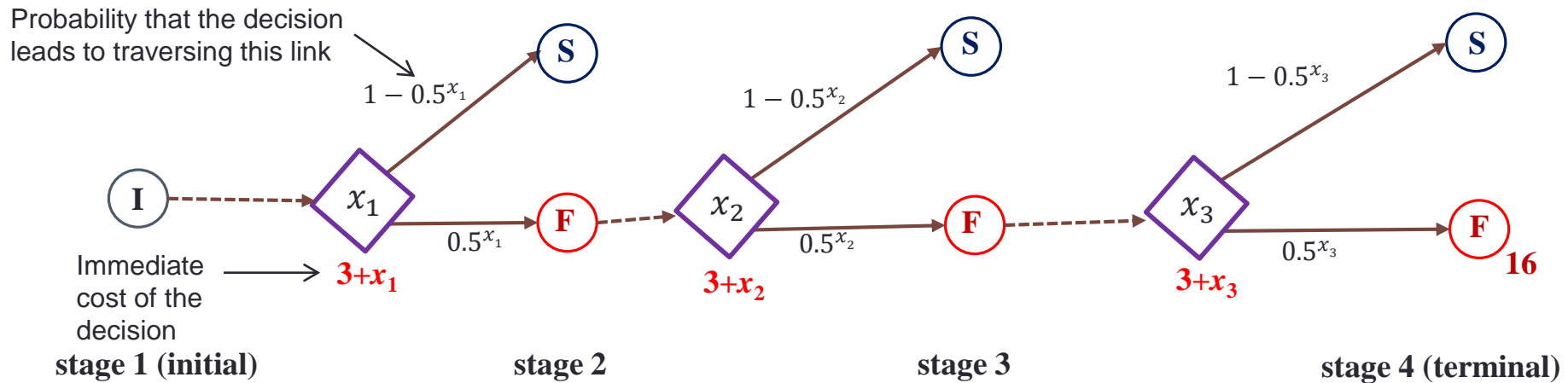
**Question: Determine the number of units to produce in each production run in order to minimize the total expected cost**

# Stages, States, Decision Vars, Optimality Criterion

- Stages: 3 stages representing decision making for each production run and a terminal stage
- States:  $S$  (success) if we do not need anymore items in stage  $i$ ,  
 $F$  (failure) if we still need an item in stage  $i$  or later
- Decision Variables:  $x_i :=$  size of the  $i$ 'th production run
- Optimality criterion: minimize total expected cost



# Value function



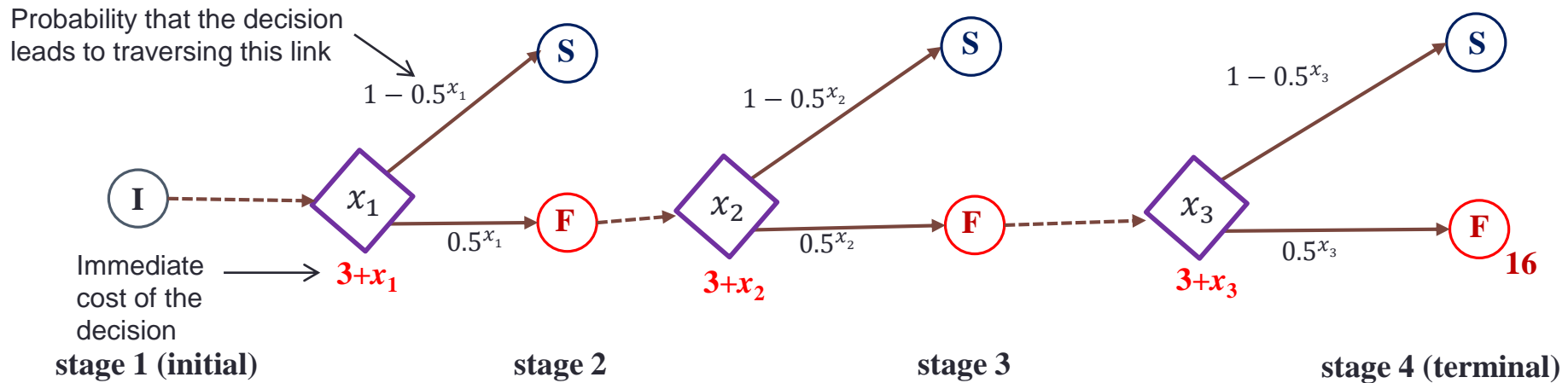
$f_n(s_n, x_n)$  = Expected cost of starting from state  $s_n$  in stage  $n$ , deciding to produce  $x_n$  and making optimal decisions thereafter

$$f_n^*(s_n) = \min_{x_n=0,1,2,3,\dots} f_n(s_n, x_n)$$

Need:  $f_1^*(I)$



# Value function: final stage



$f_n(s_n, x_n)$  = Expected cost of starting from state  $s_n$  in stage  $n$ , deciding to produce  $x_n$  and making optimal decisions thereafter

$$f_n^*(s_n) = \min_{x_n=0,1,2,3,\dots} f_n(s_n, x_n)$$

Need:  $f_1^*(I)$

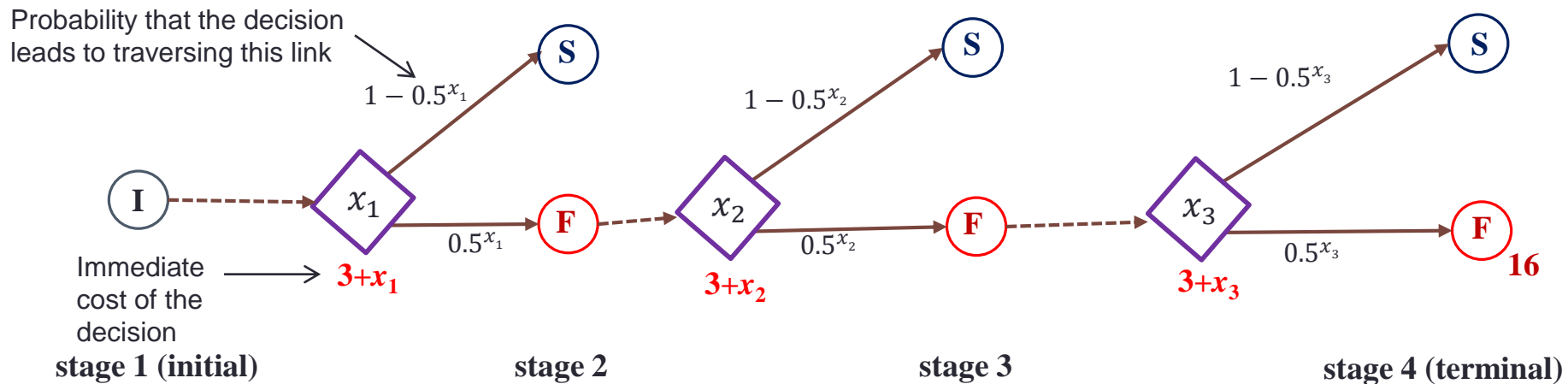
Value function if we reach success state in final stage (product accepted by the customer, no more cost)

$$f_4^*(S) = 0$$

Value function at failed state in final stage (terminal cost)

$$f_4^*(F) = 16$$

# Value function: recursive relationship



$f_n(s_n, x_n)$  = Expected cost of starting from state  $s_n$  in stage  $n$ , deciding to produce  $x_n$  and making optimal decisions thereafter

$$f_n^*(s_n) = \min_{x_n=0,1,2,3,\dots} f_n(s_n, x_n)$$

Need:  $f_1^*(I)$

Value function at a successful state in stage 2 or 3 (product accepted by the customer, no more cost)

$$f_2^*(S) = f_3^*(S) = 0$$

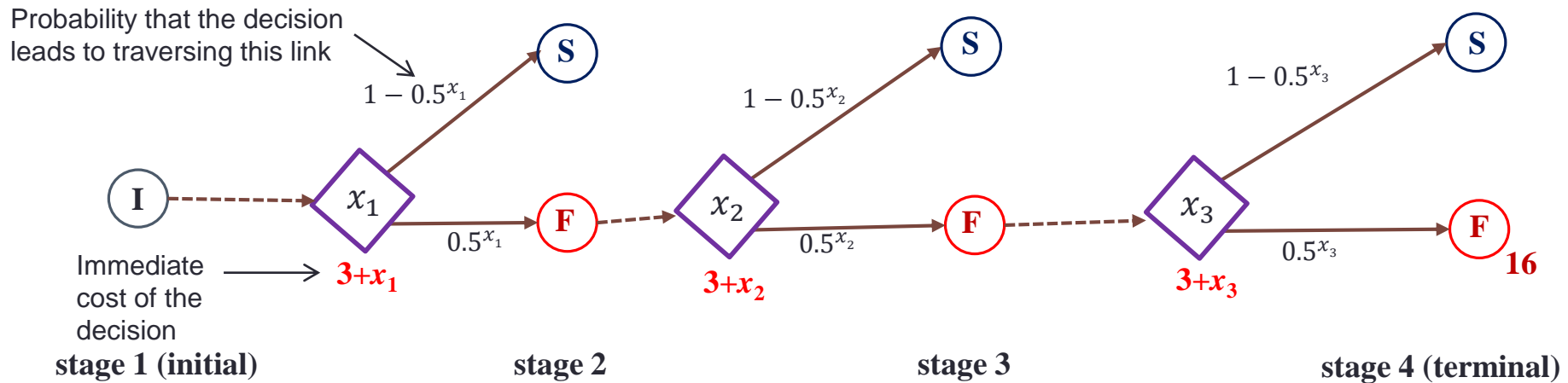
Value function at a failed state in stage 2 or 3 (expected cost of another round of production)

$$f_3(F, x_3) = 3 + x_3 + 0.5^{x_3} \times f_4^*(F)$$

$$f_3^*(F) = \min_{x_3=0,1,2,3,\dots} \{3 + x_3 + 0.5^{x_3} \times f_4^*(F)\}$$

$f_n(s_n, x_n) =$   
 Cost of producing  $x_n$  in round  $n$  (Immediate cost)  
 +  
 Probability(Success | produce  $x_n$ )  $\times$  Optimal Expected cost in the remaining stages given that we had a success  
 +  
 Probability(Failure | produce  $x_n$ )  $\times$  Optimal Expected cost in the remaining stages given that we had a failure

# Value function: recursive relationship



$f_n(s_n, x_n)$  = Expected cost of starting from state  $s_n$  in stage  $n$ , deciding to produce  $x_n$  and making optimal decisions thereafter

$$f_n^*(s_n) = \min_{x_n=0,1,2,3,\dots} f_n(s_n, x_n)$$

Need:  $f_1^*(I)$

Value function at a successful state in stage 2 or 3 (product accepted by the customer, no more cost)

$$f_2^*(S) = f_3^*(S) = 0$$

Value function at a failed state in stage 2 or 3 (expected cost of another round of production)

$$f_3^*(F) = \min_{x_3=0,1,2,3,\dots} \{3 + x_3 + 0.5^{x_3} \times f_4^*(F)\}$$

$$f_2^*(F) = \min_{x_2=0,1,2,3,\dots} \{3 + x_2 + 0.5^{x_2} \times f_3^*(F)\}$$

Value function at the initial state (similar to the value function in a failed state, need to run production)

$$f_1^*(I) = \min_{x_1=0,1,2,3,\dots} \{3 + x_1 + 0.5^{x_1} \times f_2^*(F)\}$$

## Solution Procedure: Backward induction

Need:  $f_1^*(I)$

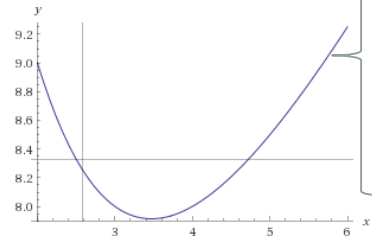
**stage 4**

$$f_4^*(F) = 16$$

**stage 3**

$$f_3^*(F) = \min_{x_3=0,1,2,3,\dots} \{3 + x_3 + 0.5^{x_3} \times f_4^*(F)\} = \min_{x_3=0,1,2,3,\dots} \{3 + x_3 + 0.5^{x_3} \times 16\}$$

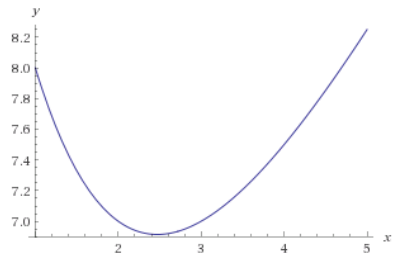
$x_3$	0	1	2	3	4	5	6	$f_3^*(F)$	$x_3^*$
$3 + x_3 + 0.5^{x_3} \times 16$	19	12	9	8	8	8.5	9.25	8	3 or 4



**stage 2**

$$f_2^*(F) = \min_{x_2=0,1,2,3,\dots} \{3 + x_2 + 0.5^{x_2} \times f_3^*(F)\} = \min_{x_2=0,1,2,3,\dots} \{3 + x_2 + 0.5^{x_2} \times 8\}$$

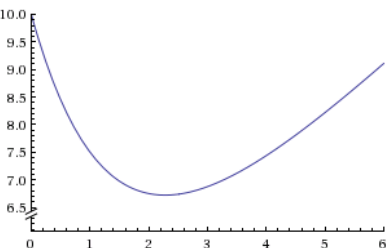
$x_2$	0	1	2	3	4	5	$f_2^*(F)$	$x_2^*$
$3 + x_2 + 0.5^{x_2} \times 8$	12	8	7	7	7.5	8.25	7	2 or 3



**stage 1**

$$f_1^*(I) = \min_{x_1=0,1,2,3,\dots} \{3 + x_1 + 0.5^{x_1} \times f_2^*(F)\} = \min_{x_1=0,1,2,3,\dots} \{3 + x_1 + 0.5^{x_1} \times 7\}$$

$x_1$	0	1	2	3	4	$f_1^*(I)$	$x_1^*$
$3 + x_1 + 0.5^{x_1} \times 7$	10	7.5	6.75	6.875	7.4375	6.75	2



## Optimal Policy

Need:  $f_1^*(I)$

Optimal expected cost is 6.75\$

An optimal policy is to produce

2 units in the first run

2 units in the second run

3 units in the third run

stage 4

$$f_4^*(F) = 16$$

stage 3

$$f_3^*(F) = \min_{x_3=0,1,2,3,\dots} \{3 + x_3 + 0.5^{x_3} \times f_4^*(F)\} = \min_{x_3=0,1,2,3,\dots} \{3 + x_3 + 0.5^{x_3} \times 16\}$$

$x_3$	0	1	2	3	4	5	6	$f_3^*(F)$	$x_3^*$
$3 + x_3 + 0.5^{x_3} \times 16$	19	12	9	8	8	8.5	9.25	8	3 or 4

stage 2

$$f_2^*(F) = \min_{x_2=0,1,2,3,\dots} \{3 + x_2 + 0.5^{x_2} \times f_3^*(F)\} = \min_{x_2=0,1,2,3,\dots} \{3 + x_2 + 0.5^{x_2} \times 8\}$$

$x_2$	0	1	2	3	4	5	$f_2^*(F)$	$x_2^*$
$3 + x_2 + 0.5^{x_2} \times 8$	12	8	7	7	7.5	8.25	7	2 or 3

stage 1

$$f_1^*(I) = \min_{x_1=0,1,2,3,\dots} \{3 + x_1 + 0.5^{x_1} \times f_2^*(F)\} = \min_{x_1=0,1,2,3,\dots} \{3 + x_1 + 0.5^{x_1} \times 7\}$$

$x_1$	0	1	2	3	4	$f_1^*(I)$	$x_1^*$
$3 + x_1 + 0.5^{x_1} \times 7$	10	7.5	6.75	6.875	7.4375	6.75	2



# WINNING IN VEGAS

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## Example 5: Winning in Vegas

- You would like to show-off your casino skills to your friend
- Friend challenges you to the following **game**:
  - You start with **3** chips and you are allowed to play **3** rounds
  - In each round,
    - you can bet either none or some, or all of the chips that you have at the beginning of that round
    - Casino odds:  $1/3$  probability of losing all chips that you bet  
 $2/3$  probability of doubling the chips that you bet
- You win the **game** if you end up with at least 5 chips after 3 rounds; otherwise you lose

**Question:** determine the number of chips to bet in each round to win the **game** against your friend